Optimizing Scheme, part I

cons should not cons its arguments, part I

a Lazv Alloc is a Smart Alloc

Alex Gal

COMP 621

cancelled

Samuel Gélineau

Optimizing Scheme, part II

an inexistant return is a smart return

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COMP 621

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stack-storage optimization for short-lived data

most object are short-lived

- allocate them on the stack (faster than malloc)
- those that outlive the function call are moved to the heap
- that's quite a short zeroth generation!

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cons should not cons its arguments. part II Chenev on the M.T.A.

Henry Baker

Sing along!

Charlie on the M.T.A.

oh. will he ever return? no, he'll never return, and his fate is still unlearned. he's a man who'll never return!



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Compiling Scheme to C Scheme and C are so different

Scheme

C

```
Hand-optimized low-level details.

void reverse(int* array, int length) {
  for(int i = 0, i = length-1; i<i; ++i, --i) {
    swap(&(array[i]), &(array[i]));
  }
}</pre>
```

No way our generated code can pull that sort of trick!

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Features only provided by C apart from segfaults

```
longimp
imp_buf handlers[MAX_DEPTH]:
int handler_depth = 0:

int trv(void (*bodv)(void)) {
   int error_code = setimp(handlers[++handler_depth]):
   if (error_code == EXIT_SUCCESS)
     bodv():
   return error_code:
}

void throw(int error_code) {
   if (error_code != EXIT_SUCCESS)
     longimp(handlers[handler_depth--]. error_code):
}
```

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Features only provided by Scheme

apart from allowing weird characters in identifiers

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a Scheme-specific optimization

required by the language definition. but not always strictly obeyed

```
void recursive_loop() {
  printf("infinite bottles of beer on the wall\n"):
  recursive_loop(): // exhausts the stack
```

tail-call optimization

Scheme

```
(define (recursive-loop)
  (display "infinite bottles of beer on the wall\n")
  (recursive-loop)) : does not exhaust the stack!

(recursive-loop)
```

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a C-specific optimization

not standard. but implemented by most compilers

*a and *c are freed at the end of the block, but not *b.

Scheme

Garbage-collection: when all vou have is a hammer...

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Amortizing the trampoline cost

"avoid making a large number of small trampoline bounces by occasionally iumping off the Empire State Building"

```
bungce
imp_buf trampoline:

void recursive_loop() {
  int _:
    printf("infinite bottles of beer on the wall\n"):
    if (&_ > STACK_LIMIT)
        longimp(trampoline, (int) recursive_loop):
    else
        recursive_loop():
}

int main() {
    bounce f = (bounce) setimp(trampoline):
    if (f == NULL) f = &recursive_loop:
    f():
}
```

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Target code for tail-recursion a bit of interpreter overhead in the compiled code

trampoline

```
trampoline

void* args;
void* result:
tvpedef void* (*bounce)():

void* recursive_loop() {
  printf("infinite bottles of beer on the wall\n"):
  return recursive_loop:
}

void trampoline() {
  bounce f = recursive_loop:
  for(:;)
    f = f();
}
```

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Garbage-collecting the stack

don't throw the live variables with the bathwater

```
a longer zeroth generation
if (&_ > STACK_LIMIT) {
  gc();
  alloca(-STACK_SIZE);
}
recursive_loop();
```

Move live variables to the heap, garbage-collect the rest.
Using a copy-collector, young dead nodes are collected for free!

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Continuation-passing-style

What if the entire program was written by a tail-call fanatic?

let all calls be tail calls

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Bungeeeeee! a one slide summarv

- never return. never.
- use continuation-passing-style to avoid returns.
- always allocate on the stack.
- when we run out of stack space:
 - flush the dead nodes (for free)
 - copy the live nodes (amortized by the mallocs we avoided)
 - flush the call stack (dec %ESP %ESP STACK_SIZE)
 - call the continuation

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